

It All Depends on How You Slice It: An Introduction to Hyperplane Arrangements

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Outline

Hyperplane Arrangements

Counting Regions

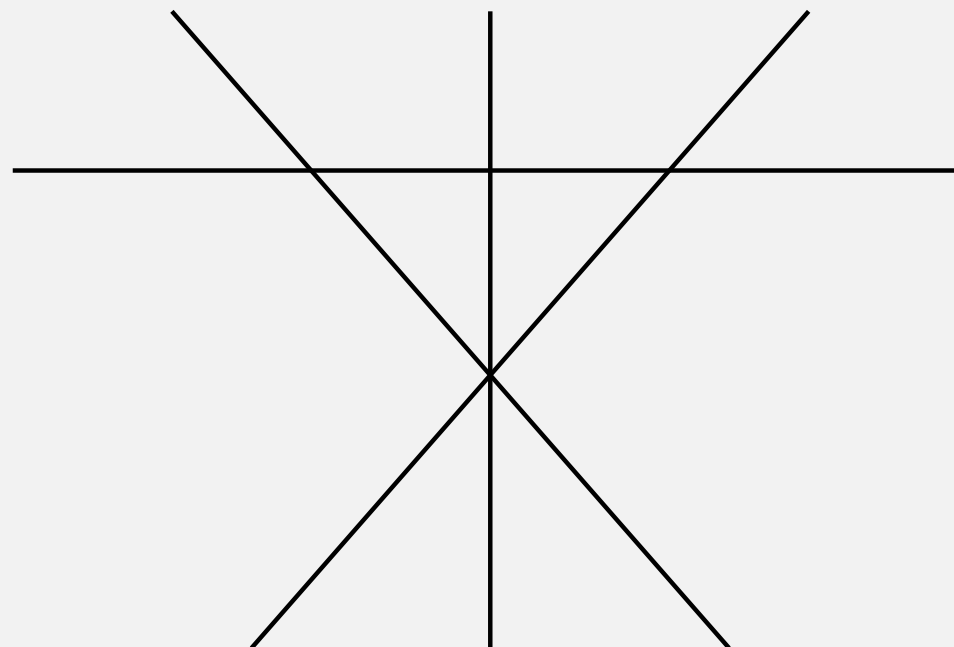
The Intersection Poset

Finite Field Method

Hyperplane Arrangements

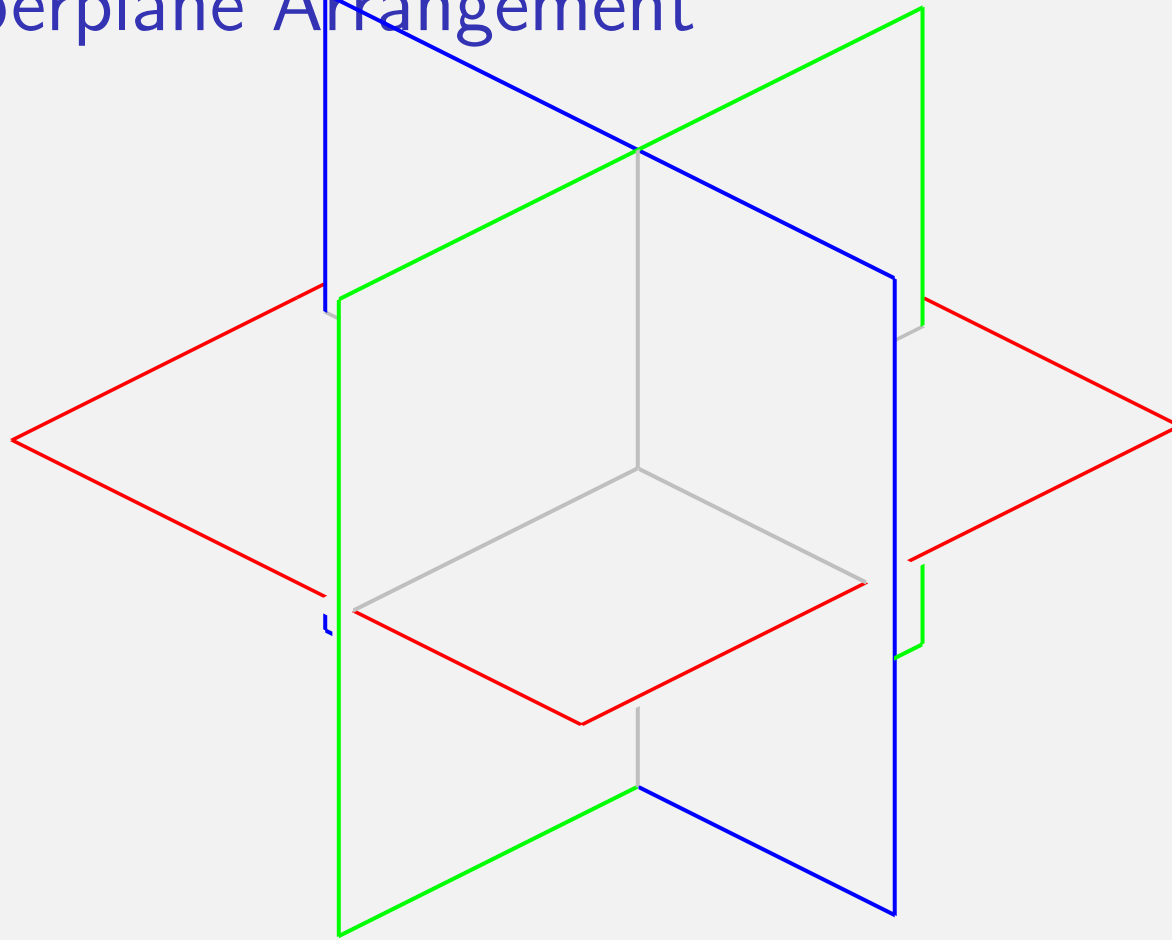
- ▶ Begin with $\mathbb{R}^n = \{(x_1, x_2, \dots, x_n) : x_i \in \mathbb{R}\}$.
- ▶ An **affine hyperplane** is the set of points in \mathbb{R}^n satisfying an equation of the form $a_1x_1 + \dots + a_nx_n = b$.
- ▶ A **hyperplane arrangement** is simply a collection of affine hyperplanes.

A Hyperplane Arrangement



An arrangement of four lines in the plane

Another Hyperplane Arrangement



An arrangement of three planes in three space

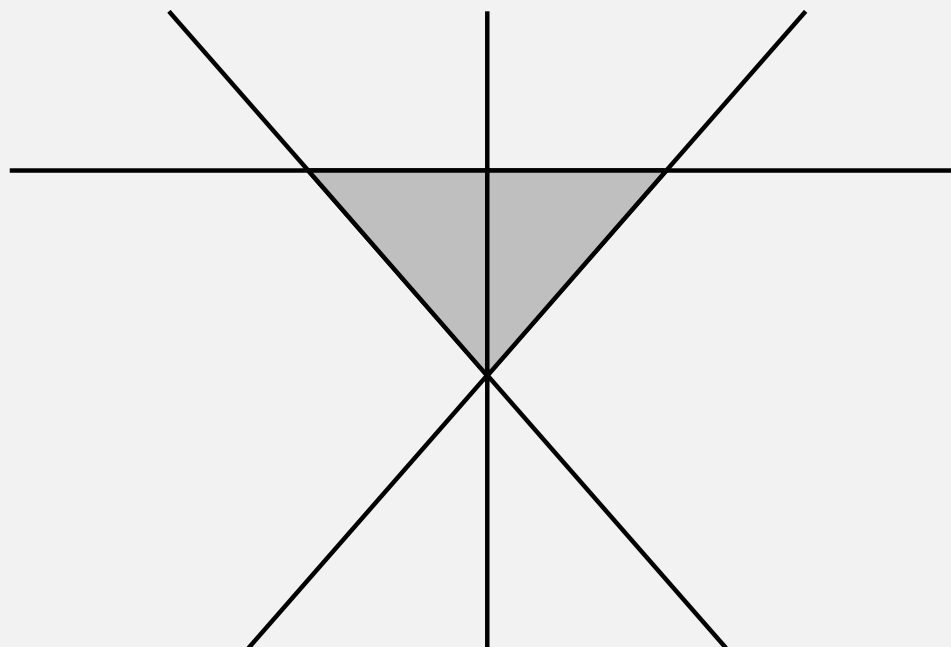
Regions

- ▶ A **region** of the arrangement \mathcal{A} is a connected component of the complement

$$\mathbb{R}^n - \bigcup_{H \in \mathcal{A}} H$$

- ▶ Regions can be bounded or unbounded.
- ▶ The total number of regions is $r(\mathcal{A})$, and the number of bounded regions is $b(\mathcal{A})$.

Region Counting



$$r(\mathcal{A}) = 10 \quad b(\mathcal{A}) = 2 \text{ (shaded)}$$

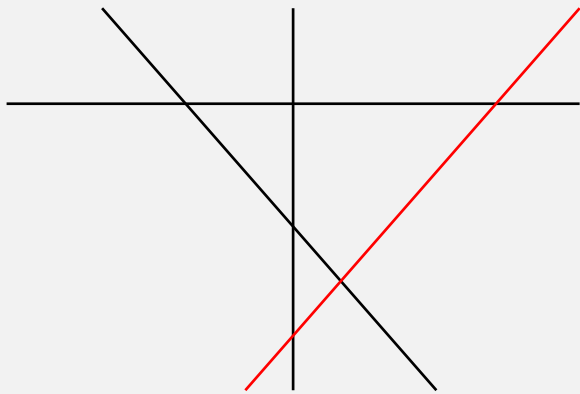
Deletion and Restriction

- ▶ Is there a better way to count regions?
- ▶ Yes!

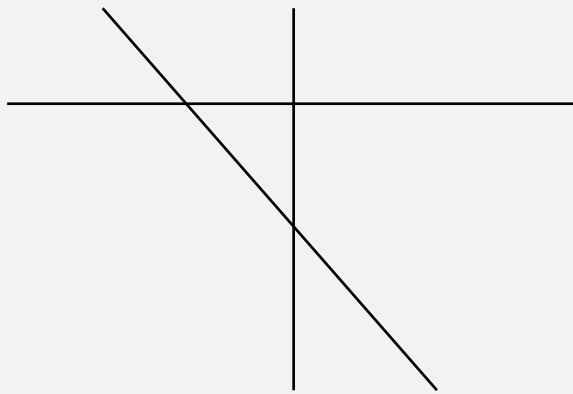
Definition

- ▶ Let \mathcal{A} be an arrangement and $H \in \mathcal{A}$ a hyperplane.
- ▶ $\mathcal{A}' = \mathcal{A} \setminus H$ is called the **deleted arrangement**.
- ▶ $\mathcal{A}'' = \{K \cap H : K \in \mathcal{A}'\}$ is called the **restricted arrangement**.
- ▶ $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ is called a **triple of arrangements**.

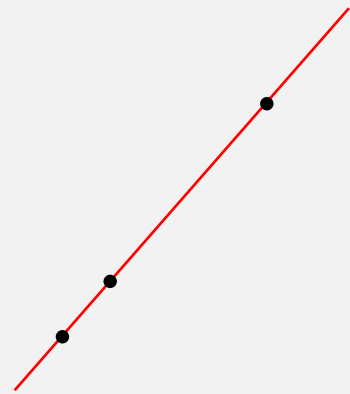
A Triple $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ of Arrangements



\mathcal{A} and H



\mathcal{A}'



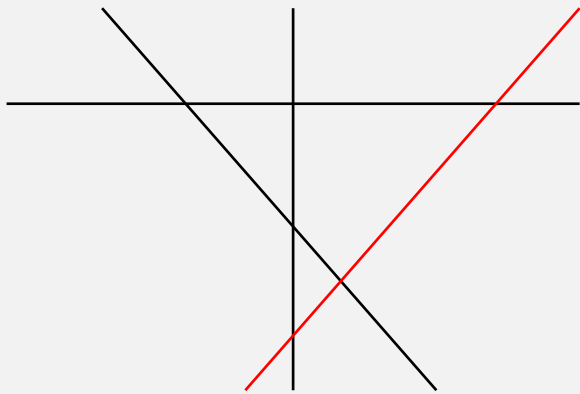
\mathcal{A}''

Region Counting Recurrence

Theorem (Zaslavsky, 1975)

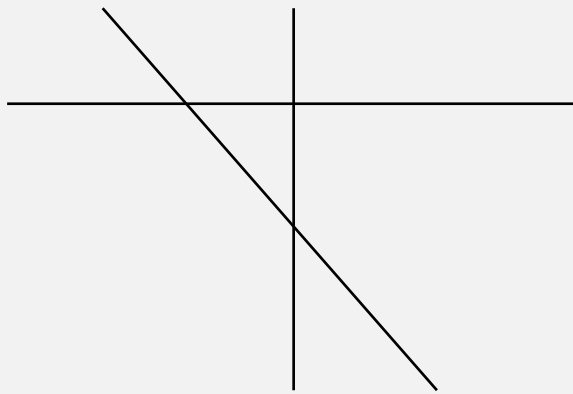
$$r(\mathcal{A}) = r(\mathcal{A}') + r(\mathcal{A}'')$$

Proof by Example (Don't try this at home!)



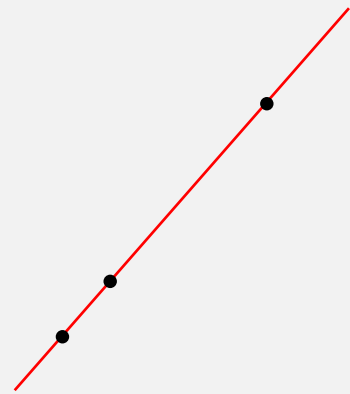
\mathcal{A} and H

$$r(\mathcal{A}) = 11$$



\mathcal{A}'

$$r(\mathcal{A}') = 7$$



\mathcal{A}''

$$r(\mathcal{A}'') = 4$$

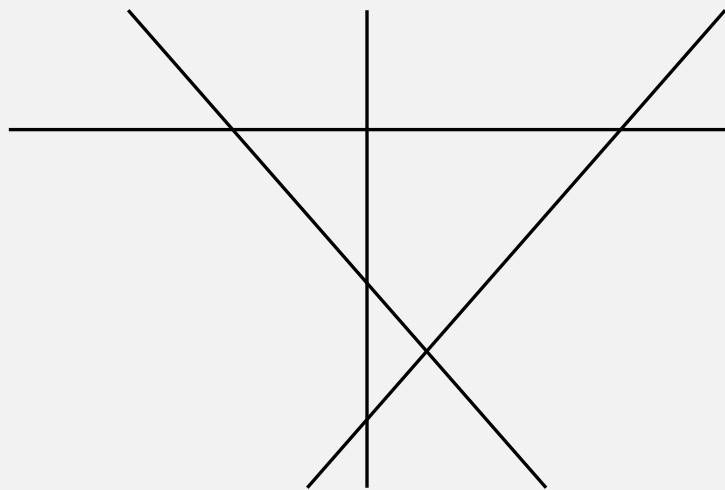
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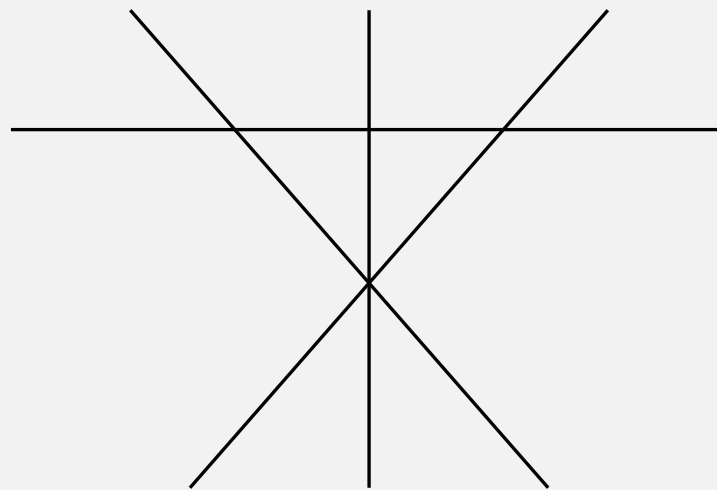
Arrangements in General Position

- ▶ Let's apply this result.
- ▶ An arrangement \mathcal{A} is in **general position** if you can move the hyperplanes slightly and not change the number of regions.

Two Arrangements



in general position



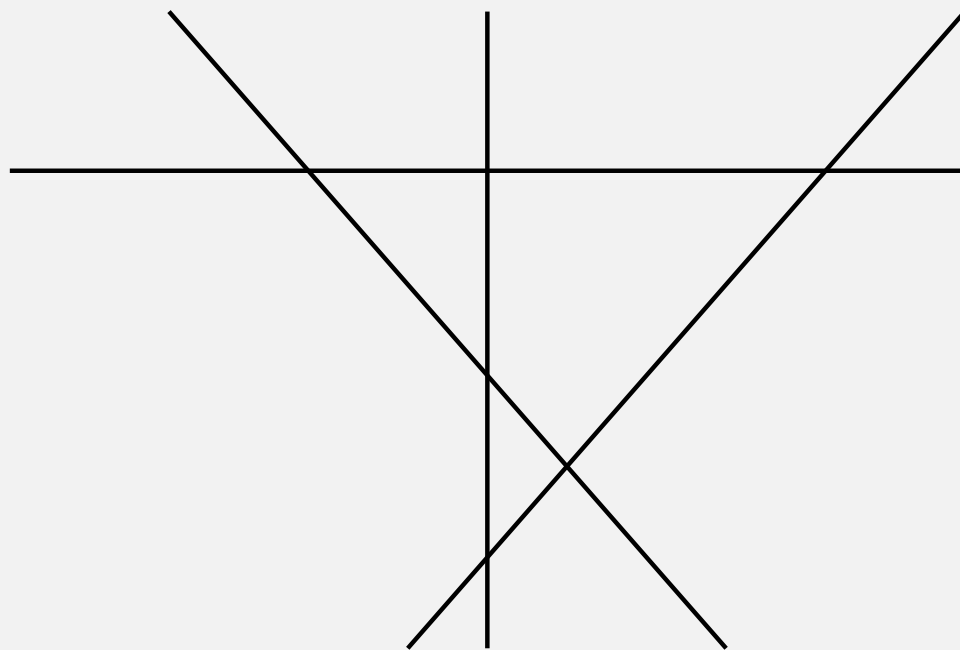
not in general position

Counting Regions of General Position Line Arrangements

- ▶ Start with an arrangement \mathcal{A} of k lines in general position in the plane, and choose a particular line H .
- ▶ By hypothesis, H meets \mathcal{A}' in $k - 1$ points, which divide H into k regions. So $r(\mathcal{A}'') = k$.
- ▶ Hence $r(\mathcal{A}) = r(\mathcal{A}') + k$, where \mathcal{A}' contains $k - 1$ lines.
- ▶ By continuing to delete lines in this way, we get $r(\mathcal{A}) = r(\emptyset) + 1 + 2 + \cdots + (k - 1) + k$.
- ▶ When no lines are present there is one region, so $r(\emptyset) = 1$.
- ▶ Hence, for a general position line arrangement we have

$$r(\mathcal{A}) = 1 + \sum_{i=1}^k i = 1 + \frac{k(k+1)}{2} = 1 + k + \binom{k}{2}.$$

Counting Regions of a General Position Line Arrangement



$$r(\mathcal{A}) = 1 + k + \binom{k}{2} = 1 + 4 + \binom{4}{2} = 11 \checkmark$$

The Number of Regions in an Arbitrary Arrangement

- ▶ We can use the recurrence (and induction) to show that

$$r(\mathcal{A}) = 1 + k + \binom{k}{2} + \binom{k}{3} + \cdots + \binom{k}{n}$$

for k hyperplanes in general position in n dimensional space.
(Ludwig Schläfli, 1901)

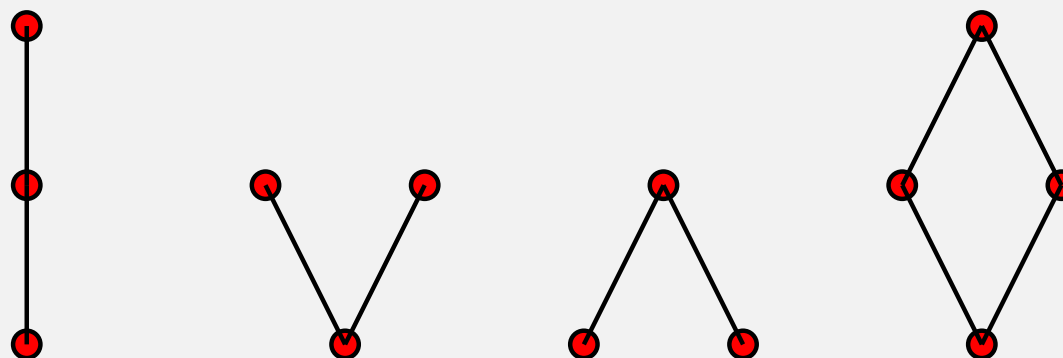
- ▶ But what if the hyperplanes are not in general position?
- ▶ Zaslavsky developed a powerful tool to compute $r(\mathcal{A})$ in general. To describe this, we must take a long detour...

Partially Ordered Sets

A **partially ordered set (poset)** is a set P and a relation \leq satisfying the following axioms (for all x , y , and z in P):

1. (reflexivity) $x \leq x$.
2. (antisymmetry) $x \leq y$ and $y \leq x$ implies $x = y$.
3. (transitivity) $x \leq y$ and $y \leq z$ implies $x \leq z$.
4. Posets are represented by their (Hasse) diagrams.

Hasse Diagrams

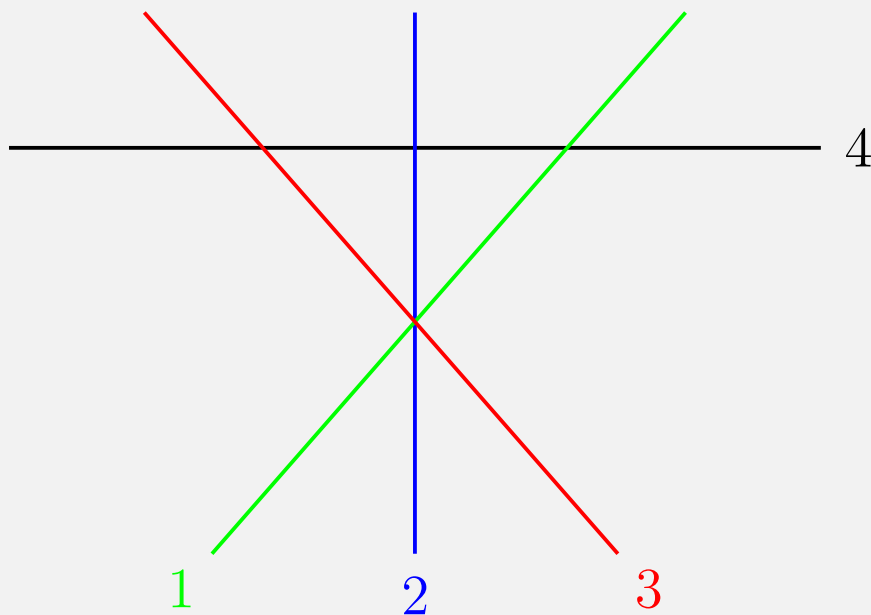


Some posets

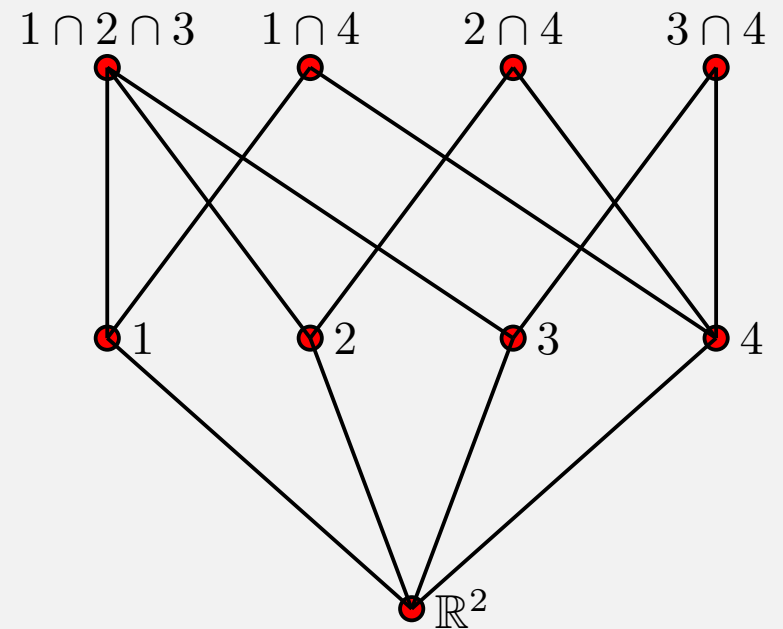
The Intersection Poset of an Arrangement

- ▶ The **intersection poset** $L(\mathcal{A})$ of the arrangement \mathcal{A} has as its elements all the intersections of all the hyperplanes.
- ▶ It is ordered (for good reason) by **reverse inclusion**, so $A \subseteq B \Leftrightarrow A \geq B$.
- ▶ The minimum element is the ambient space \mathbb{R}^n .

An Arrangement and Its Intersection Poset



A labeled arrangement \mathcal{A}



Its intersection poset $L(\mathcal{A})$

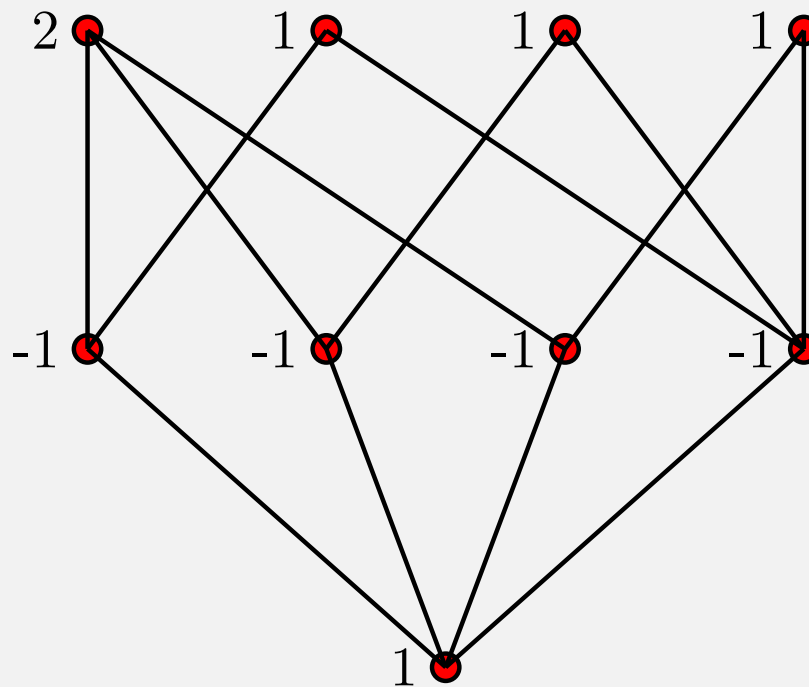
The Möbius Function for Posets

- ▶ The (closed) **interval** $[x, y]$ of a poset is the set of all points between x and y , including endpoints:

$$[x, y] = \{z : x \leq z \leq y\}.$$

- ▶ The **Möbius function** $\mu(x, y)$ is defined (recursively) on the interval $[x, y]$ by the two properties:
 1. $\mu(x, x) = 1$.
 2. $x < y \Rightarrow \sum_{z \in [x, y]} \mu(x, z) = 0$

Some Möbius Function Values



The values of the Möbius functions $\mu(\hat{0}, x)$

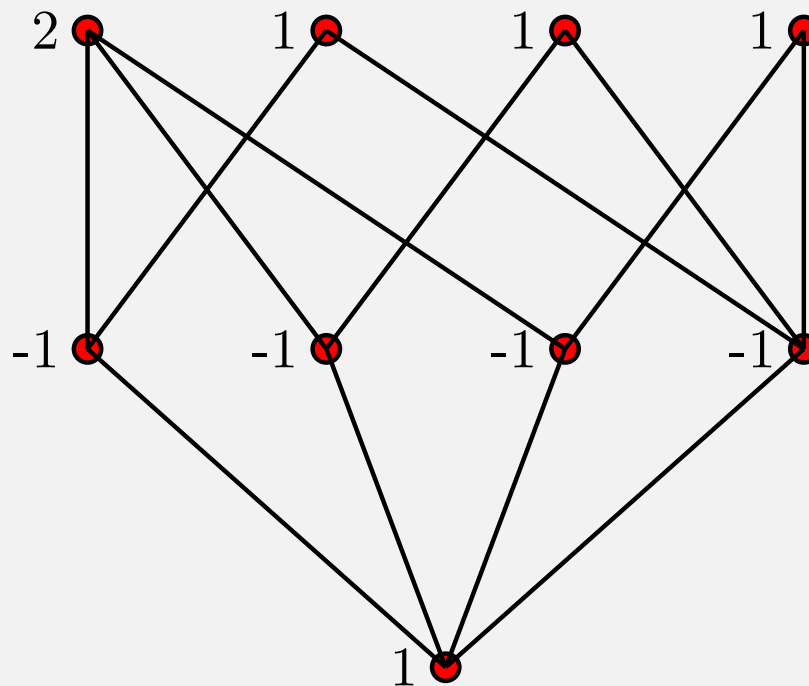
The Characteristic Polynomial

- ▶ We define the **characteristic polynomial** associated to the arrangement \mathcal{A} by

$$\chi(\mathcal{A}, q) = \sum_{x \in L(\mathcal{A})} \mu(\hat{0}, x) q^{\dim(x)}$$

.

The Characteristic Polynomial of an Arrangement



$$\chi(\mathcal{A}, q) = \sum_{x \in L(\mathcal{A})} \mu(\hat{0}, x) q^{\dim(x)} = q^2 - 4q + 5.$$

Zaslavsky's Theorem

Theorem (Zaslavsky, 1975)

With the definitions above

$$r(\mathcal{A}) = \sum_{x \in L(\mathcal{A})} |\mu(\hat{0}, x)| = |\chi(\mathcal{A}, -1)|$$

$$b(\mathcal{A}) = \left| \sum_{x \in L(\mathcal{A})} \mu(\hat{0}, x) \right| = |\chi(\mathcal{A}, 1)|$$

Zaslavsky's Theorem—continued

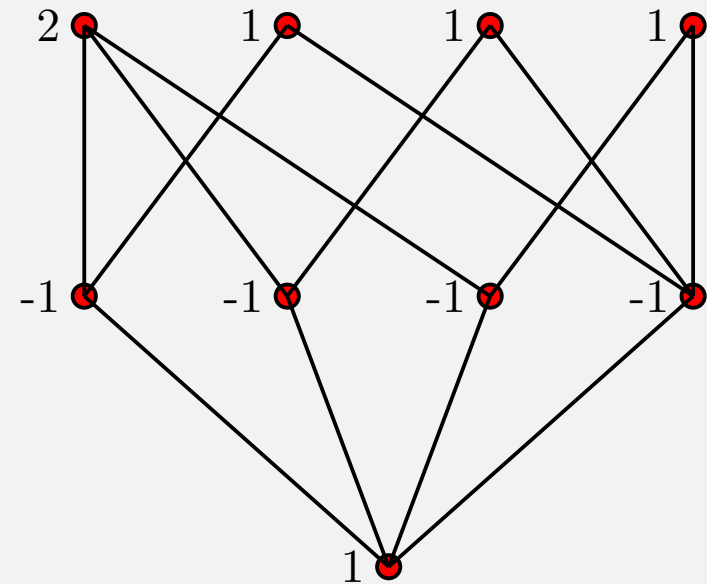
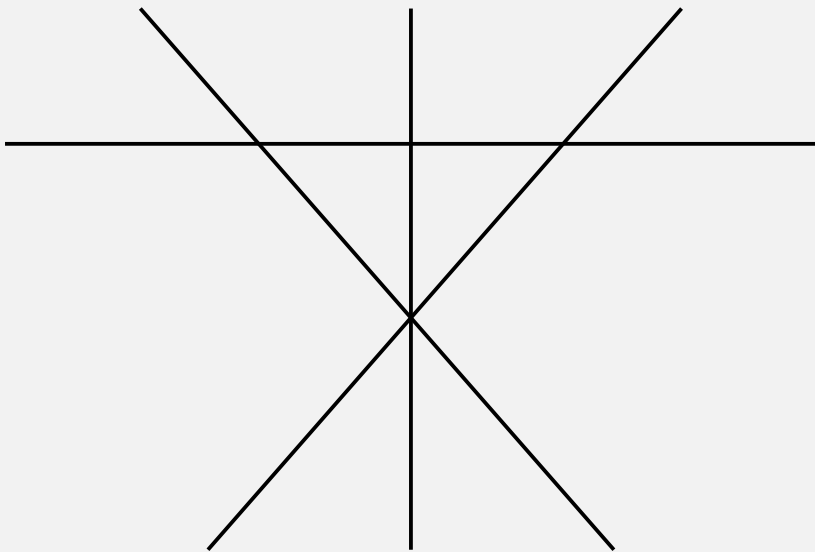
Idea of Proof.

One can show that, for any triple $(\mathcal{A}, \mathcal{A}', \mathcal{A}'')$ of arrangements,

$$\chi(\mathcal{A}, q) = \chi(\mathcal{A}', q) - \chi(\mathcal{A}'', q),$$

from which it follows that $|\chi(\mathcal{A}, -1)|$ and $r(\mathcal{A})$ satisfy the same recurrence. As they agree on the empty set, the claim follows. A similar argument works for $b(\mathcal{A})$. □

Counting Regions *via* Zaslavsky's Theorem



$$\chi(\mathcal{A}, q) = q^2 - 4q + 5$$

$$r(\mathcal{A}) = |\chi(\mathcal{A}, -1)| = 10$$

$$b(\mathcal{A}) = |\chi(\mathcal{A}, 1)| = 2$$

Finite Field Method

- ▶ Recall that the defining equation of a hyperplane can be written $a_1x_1 + \cdots + a_nx_n = b$ for some real numbers $\{a_1, a_2, \dots, a_n, b\}$.
- ▶ In many cases of interest the numbers a_1, a_2, \dots, a_n, b are integers.
- ▶ When this holds there is a particularly nice way to compute the characteristic polynomial.
- ▶ For any positive integer q let \mathcal{A}_q denote the hyperplane arrangement \mathcal{A} with defining equations reduced mod q .
- ▶ Then we have the following result.

The Characteristic Polynomial for Integral Arrangements

Theorem (Crapo and Rota (1971), Orlik and Terao (1992), Athanasiadis (1996), Björner and Ekedahl (1996))

For q a sufficiently large prime

$$\chi(\mathcal{A}, q) = \# \left(\mathbb{F}_q^n - \bigcup_{H \in \mathcal{A}_q} H \right)$$

where \mathbb{F}_q^n denotes the vector space of dimension n over the finite field with q elements.

Remarks

- ▶ Identifying \mathbb{F}_q^n with $\{0, 1, \dots, q - 1\}^n = [0, q - 1]^n$, $\chi(\mathcal{A}, q)$ is the number of points in $[0, q - 1]^n$ that do *not* satisfy modulo q the defining equations of any of the hyperplanes in \mathcal{A} .
- ▶ We need large q to avoid lowering the rank of the defining matrix, but as both sides are polynomials in q , the two sides will agree for all q .

Reflection Arrangements

- ▶ Consider the following families of arrangements:

$$\mathcal{A}_n = \{x_i - x_j = 0 \mid 1 \leq i < j \leq n\}$$

$$\mathcal{D}_n = \mathcal{A}_n \cup \{x_i + x_j = 0 \mid 1 \leq i < j \leq n\}$$

$$\mathcal{B}_n = \mathcal{D}_n \cup \{x_i = 0 \mid 1 \leq i \leq n\}$$

- ▶ These are examples of reflection arrangements associated to finite Coxeter groups of types A_{n-1} , D_n , and B_n , respectively.

Computing the Characteristic Polynomial I

- ▶ What is $\chi(\mathcal{A}_n)$?
- ▶ According to the finite field method, we want the number of points in $[0, q - 1]^n$ satisfying $x_i \neq x_j$ for all $1 \leq i < j \leq n$.
- ▶ This is the same thing as asking for vectors (x_1, x_2, \dots, x_n) all of whose entries are distinct mod q .
- ▶ Well, we can pick x_1 in q ways, then x_2 in $q - 1$ ways, and so on. Thus

$$\chi(\mathcal{A}_n, q) = q(q - 1)(q - 2) \cdots (q - n + 1).$$

- ▶ It follows that $r(\mathcal{A}_n) = n!$ (and $b(\mathcal{A}_n) = 0$).

Computing the Characteristic Polynomial II

- ▶ What is $\chi(\mathcal{B}_n)$?
- ▶ Now we want to count the points satisfying $x_i \neq x_j$, $x_i \neq -x_j$, and $x_i \neq 0$.
- ▶ Since we do not allow 0, there are only $q - 1$ (nonzero) choices for the first entry, $q - 3$ nonzero choices for the second entry (because we must avoid the first entry and its negative), *etc.*.
- ▶ Thus

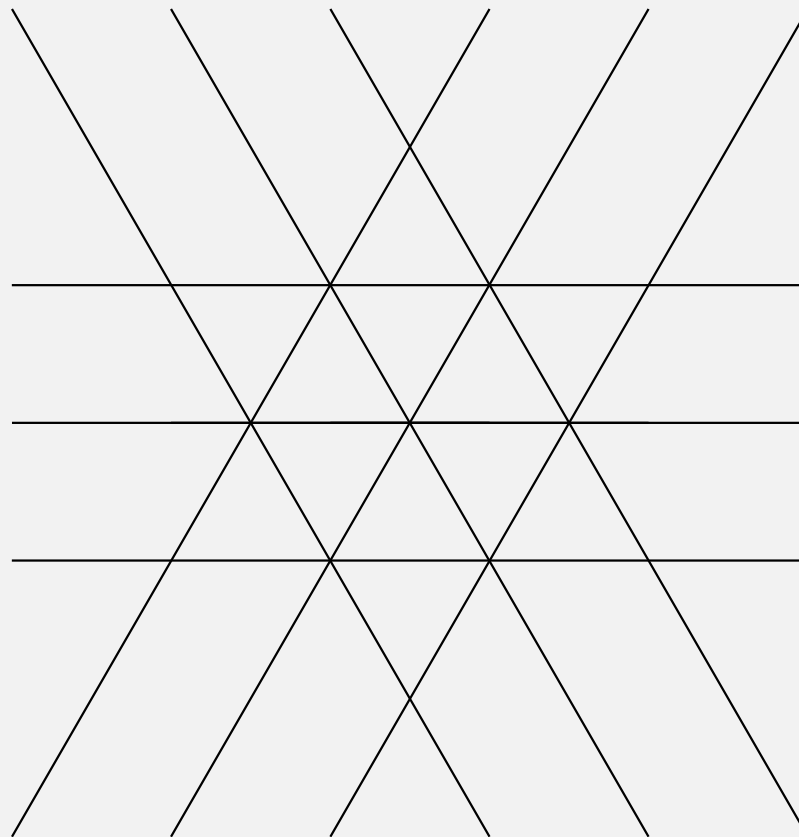
$$\chi(\mathcal{B}_n, q) = (q - 1)(q - 3) \cdots (q - 2n + 1).$$

- ▶ It follows that $r(\mathcal{B}_n) = 2^n n!$ (and $b(\mathcal{B}_n) = 0$).

The Catalan Arrangement

$$\mathcal{C}_n = \{x_i - x_j = -1, 0, +1 \mid 1 \leq i \leq j \leq n\}$$

\mathcal{C}_3 (projected)



The Catalan Arrangement

- ▶ Using the finite field method you can show that

$$\chi(\mathcal{C}_n, q) = q(q - n - 1)(q - n - 2) \cdots (q - 2n + 1)$$

- ▶ Hence

$$r(\mathcal{C}_n) = n!C_n \quad \text{and} \quad b(\mathcal{C}_n) = n!C_{n-1}$$

where

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

is the n^{th} **Catalan number**.

- ▶ As of April 1, 2008, Richard Stanley has listed **164** combinatorial interpretations of C_n on his website.

Fundamental Open Question

When does the characteristic polynomial factor completely over the integers?